



Rainbow Tables

Yukai Zang

Division of Science and Mathematics

University of Minnesota, Morris

Morris, Minnesota, USA 56267

zangx040@morris.umn.edu



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- Rainbow table
- Create rainbow tables (offline stage)
- Use rainbow tables (online stage)
- Tests
- Conclusion



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Introduction & Background



Introduction & Background

Your password
(plain-text)

Fox

Hash function

Hashed value

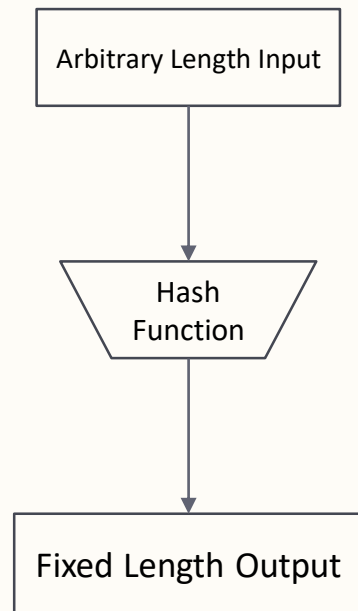
4C5E 9S8D D8S9
5T8V A7SE ASD9

Data
base



Introduction & Background

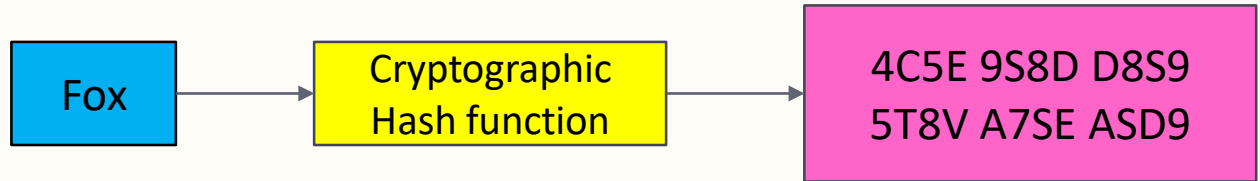
- Hash function
 - Map data of arbitrary size onto data of fixed size





Introduction & Background

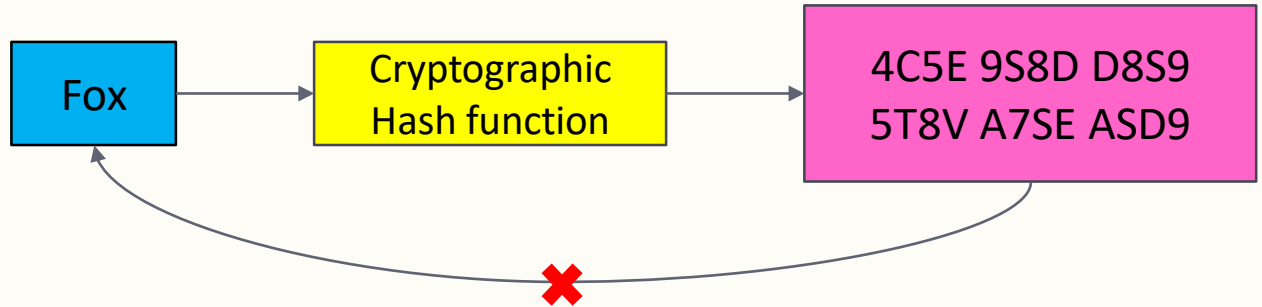
- Cryptographic hash function
 - Same plain-text result in same hashed value;





Introduction & Background

- Cryptographic hash function
 - Same plain-text result in same hashed value;
 - Fast to compute;
 - Infeasible to revert back to plain-text from hashed value;



A decorative graphic of a feather, rendered in a light beige or tan color, positioned on the left side of the slide. The feather has a central rachis with numerous barbs extending outwards, creating a fan-like shape. It is oriented vertically, pointing downwards.

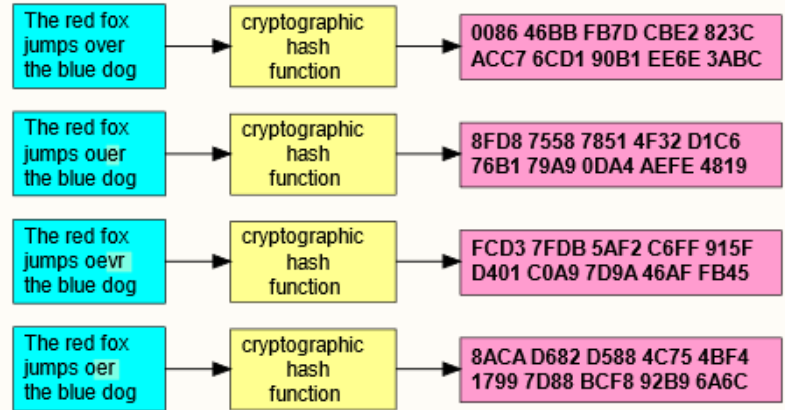
Introduction & Background

- Cryptographic hash function
 - Same plain-text result in same hashed value;
 - Fast to compute;
 - Infeasible to revert back to plain-text from hashed value;
 - Small change(s) in plain-text will cause huge changes in hashed value;

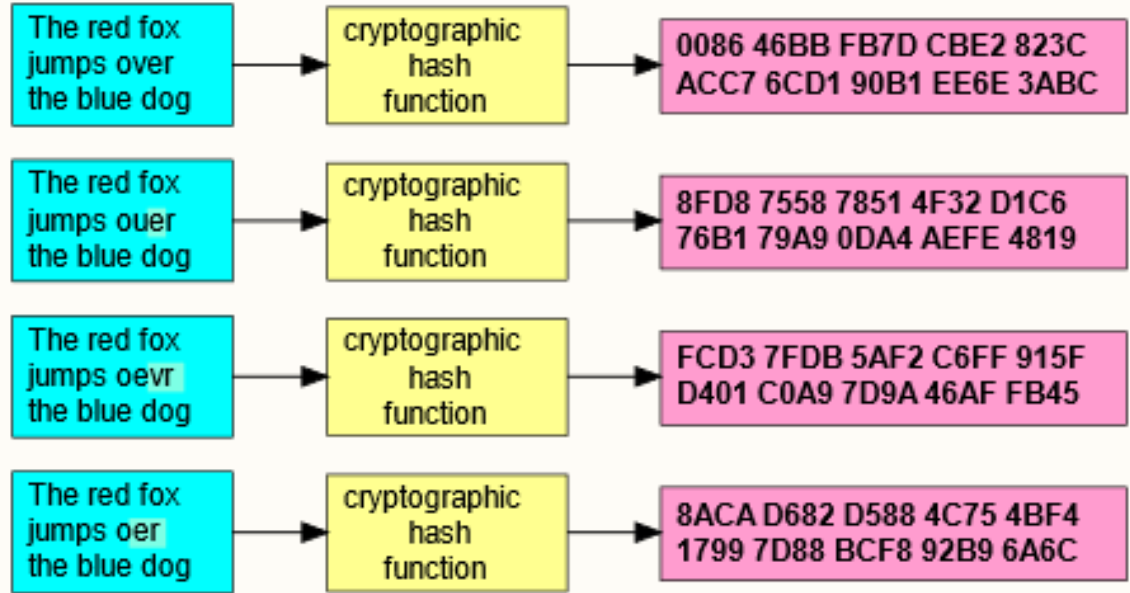


Introduction & Background

- Cryptographic hash function
 - Small change(s) in plain-text will cause huge changes in hashed value;



Introduction & Background





Introduction & Background

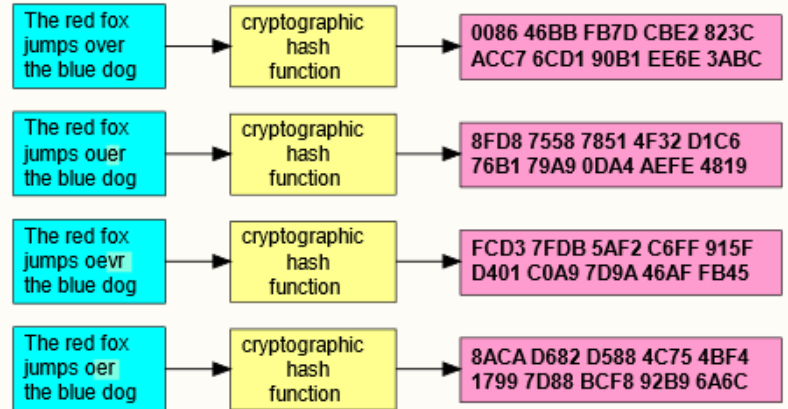
- Cryptographic hash function
 - Same plain-text result in same hashed value;
 - Fast to compute;
 - Infeasible to revert back to plain-text from hashed value;
 - Small change(s) in plain-text will cause huge changes in hashed value;
 - Infeasible to find two different plain-text with the same hashed value.



Introduction & Background

- Cryptographic hash function
 - Same plain-text result in same hashed value;

- Brute force



Introduction & Background

Password length	Alphanumeric	Days
1	62	0
2	3844	0
3	238328	0
4	14776336	0
5	916132832	0
6	56800235584	0
7	3.52 e+12	1
8	2.18 e+12	42
9	1.35 e+12	2599
10	8.39 e+12	161156

- Lookup table
- 8 digits; 62 characters
- 222 trillion combinations
- 2 quadrillion bytes
= 1800 terabytes
- 32 days with i7

Plain-text	Hash value
aaaaaaaa	7D9SXF
aaaaaaab	7WS4G5
aaaaaaac	5F2V6D
...	...
00000000	8CVIDF
00000001	1QSD9F
...	...



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- **Rainbow table**
- Create rainbow tables (offline stage)
- Use rainbow tables (online stage)
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Rainbow table

- Based on the idea introduced by Martin Hellman in 1980
- Improved by Philippe Oechslin in 2003
- Two stages
 - Offline stage
 - Online stage



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Rainbow table

- Offline stage
 - Create tables
 - *Chains*
 - *Cryptographic hash function (h)*
 - *Reduction function (r)*
 - *Collision*



Rainbow table

- Chains
 - Only start point t_0 and end point h_2 matters
- Tables look like:

$$S_0 = t_0 \xrightarrow{h} h_0 \xrightarrow{r_0} t_1 \xrightarrow{h} h_1 \xrightarrow{r_1} t_2 \xrightarrow{h} h_2 = E_0$$

Text		Hash
aaaa	↔	3c62c99bafbe3a81f16d36b592b66ae5
aaab	↔	7991ef7c61d85c505adb63b70553b0ce
aaac	↔	ef834a79dc0ac78e355859ec4f698d0b
aaad	↔	3bb75d8a8effe7c5e18aC6670a3c3346
⋮	↔	⋮



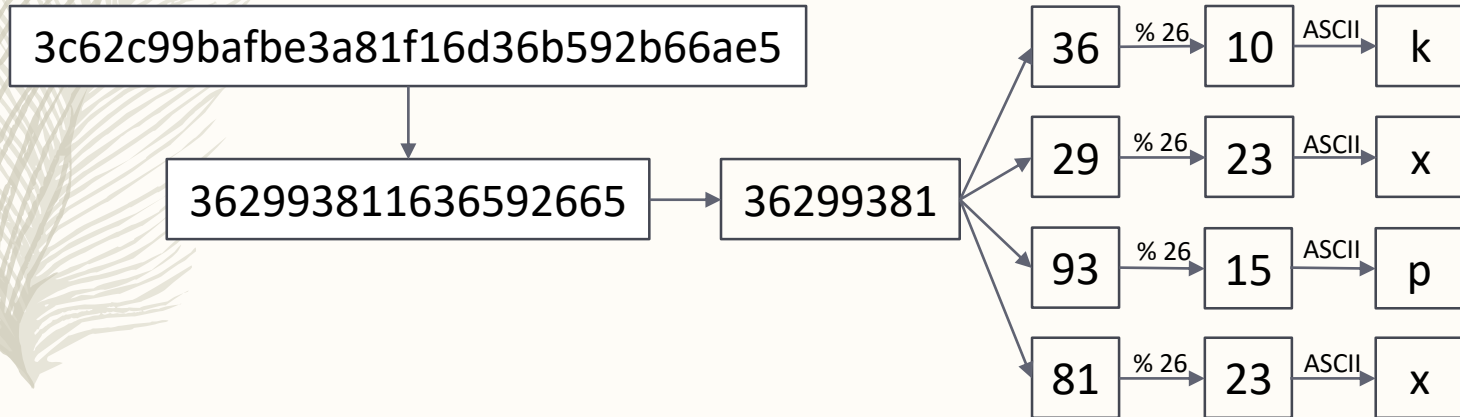


Rainbow table

- Cryptographic hash function (h)
 - MD5, SHA-1, SHA-2, SHA-3, BLAKE2, and etc.
 - Slight different in time due to different algorithm used

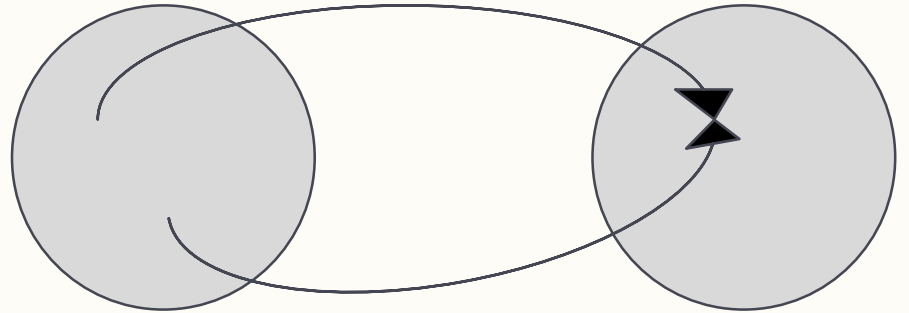
Rainbow table

- Reduction function (r)
 - Map data from the set of hashed values to the set of all plain-texts
 - Example:



Rainbow table

- Collision
 - When we map two different thing into the same value, no matter which way, it is called collision



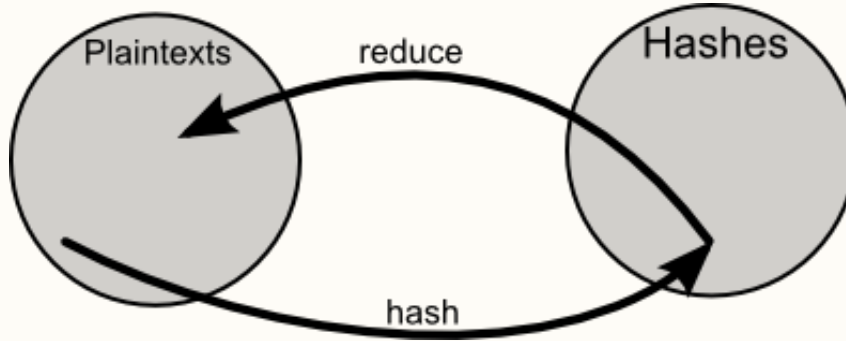


Rainbow table

- Collision
 - The size of the set of plain-texts and the set of hashed value will be different normally
 - When we try to map from a larger set onto a smaller set, collisions occur more frequently (Pigeonhole principle)

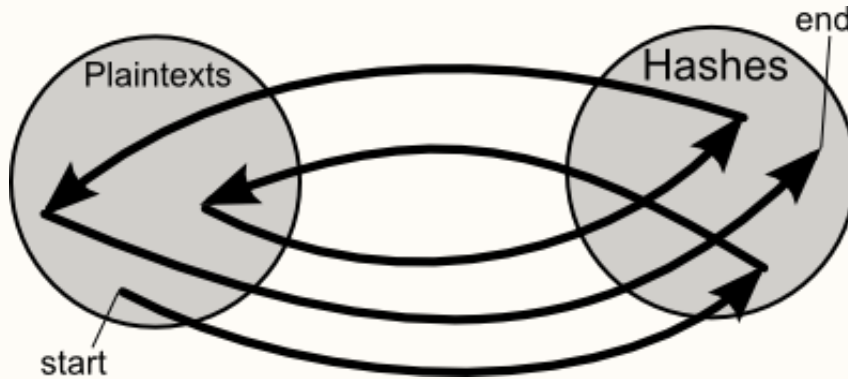
Rainbow table

- Generate steps:



Rainbow table

- Generate steps:



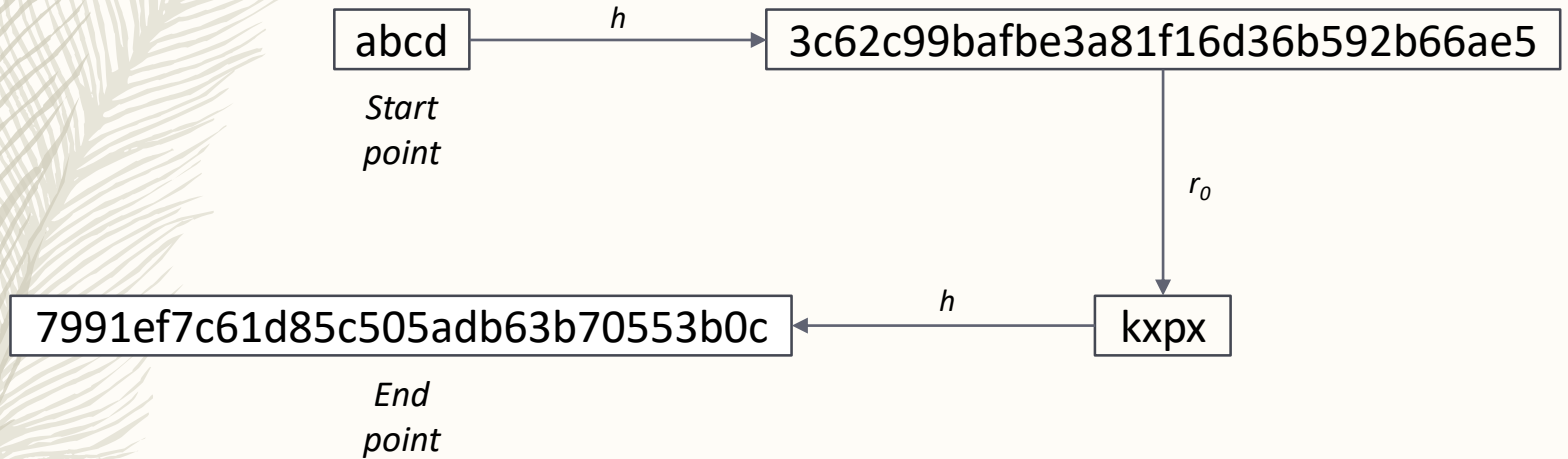
Rainbow table

- The structure of a rainbow table:

$S_0 = t_{0,0}$	\xrightarrow{h}	$h_{0,0}$	$\xrightarrow{r_0 \circ h}$	$h_{0,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{0,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{0,n} = E_0$
$S_1 = t_{1,0}$	\xrightarrow{h}	$h_{1,0}$	$\xrightarrow{r_0 \circ h}$	$h_{1,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{1,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{1,n} = E_1$
\vdots		\vdots		\vdots		\vdots		\vdots		\vdots
$S_j = t_{j,0}$	\xrightarrow{h}	$h_{j,0}$	$\xrightarrow{r_0 \circ h}$	$h_{j,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{j,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{j,n} = E_j$
\vdots		\vdots		\vdots		\vdots		\vdots		\vdots
$S_m = t_{m,0}$	\xrightarrow{h}	$h_{m,0}$	$\xrightarrow{r_0 \circ h}$	$h_{m,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{m,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{m,n} = E_m$

Rainbow table

– Example:





Rainbow table

- Example (collision):

$$t_0 = \text{"abcd"} \rightarrow r_0(h(t_0)) = \text{"defg"}$$

$$t_1 = \text{"abdc"} \rightarrow r_0(h(t_1)) = \text{"defg"}$$

we will have two chains with different start points but the same end point.

- Clean table
 - Only keep one chain with the same end point.



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Rainbow table

- Online stage
 - Actual search
 - Example:

We have a rainbow table with chain length = 3:

Rainbow table: $t_0 \sim \rightarrow h_2$

Chain: $t_0 \xrightarrow{h} h_0 \xrightarrow{r_0} t_1 \xrightarrow{h} h_1 \xrightarrow{r_1} t_2 \xrightarrow{h} h_2$

Assuming the password we want to

search for hash value:

3C62C9 & 77CC7F

Text	Hash
aaaa	3C62C9
aaab	7991EF
aaac	EF834A
spkn	3BB75D

$S_0 = t_{0,0}$	\xrightarrow{h}	$h_{0,0}$	$\xrightarrow{r_0 \circ h}$	$h_{0,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{0,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{0,n} = E_0$
$S_1 = t_{1,0}$	\xrightarrow{h}	$h_{1,0}$	$\xrightarrow{r_0 \circ h}$	$h_{1,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{1,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{1,n} = E_1$
\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
$S_j = t_{j,0}$	\xrightarrow{h}	$h_{j,0}$	$\xrightarrow{r_0 \circ h}$	$h_{j,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{j,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{j,n} = E_j$
\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
$S_m = t_{m,0}$	\xrightarrow{h}	$h_{m,0}$	$\xrightarrow{r_0 \circ h}$	$h_{m,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{m,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{m,n} = E_m$

Chain: $t_0 \xrightarrow{h} h_0 \xrightarrow{r_0} t_1 \xrightarrow{h} h_1 \xrightarrow{r_1} t_2 \xrightarrow{h} h_2$

Text	Hash
aaaa	3C62C9
aaab	7991EF
aaac	EF834A
spkn	3BB75D

1. According to the given hash value $y=h(x)=3C62C9$, search through the column for the ending points in all tables;
2. We find it at the first row with the plain-text “aaaa”;
3. We create a chain from “aaaa”, and search for the original hash (3C62C9) in the chain; if we find it, it returns the password at a given index (indexHash -1). The chain is:

$aaaa \xrightarrow{h} 584C19 \xrightarrow{r_0} kcyI \xrightarrow{h} 48950E$
 $\xrightarrow{r_1} \mathbf{puwr} \xrightarrow{h} 3C62C9$

4. In this case, the plain-text we wanted is “puwr”.

$S_0 = t_{0,0}$	\xrightarrow{h}	$h_{0,0}$	$\xrightarrow{r_0 \circ h}$	$h_{0,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{0,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{0,n} = E_0$
$S_1 = t_{1,0}$	\xrightarrow{h}	$h_{1,0}$	$\xrightarrow{r_0 \circ h}$	$h_{1,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{1,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{1,n} = E_1$
\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
$S_j = t_{j,0}$	\xrightarrow{h}	$h_{j,0}$	$\xrightarrow{r_0 \circ h}$	$h_{j,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{j,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{j,n} = E_j$
\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots
$S_m = t_{m,0}$	\xrightarrow{h}	$h_{m,0}$	$\xrightarrow{r_0 \circ h}$	$h_{m,1}$	$\xrightarrow{r_1 \circ h}$	\dots	$\xrightarrow{r_{n-2} \circ h}$	$h_{m,n-1}$	$\xrightarrow{r_{n-1} \circ h}$	$h_{m,n} = E_m$

Chain: $t_0 \xrightarrow{h} h_0 \xrightarrow{r_0} t_1 \xrightarrow{h} h_1 \xrightarrow{r_1} t_2 \xrightarrow{h} h_2$

Text	Hash
aaaa	3C62C9
aaab	7991EF
aaac	EF834A
spkn	3BB75D

1. According to the given hash value $y=h(x)=77CC7F$, search through the column for the ending points in all tables;
2. There is no matched E_j .
3. Compute the $h(r_1(77CC7F)) = 48950E$;
4. $48950E$ is not in the table;
5. Compute the $h(r_1(h(r_0(77CC7F))) = 3BB75D$;
6. We have a matched E_j , which is the last row with plain-text “spkn”;
7. We create a chain from “spkn”, and search for the original hash ($77CC7F$) in the chain; if we find it, it returns the password at a given index (indexHash -1). The chain is:

$$\text{spkn} \xrightarrow{h} 77CC7F \xrightarrow{r_0} \text{puwr} \xrightarrow{h} 584C19$$

$$\xrightarrow{r_1} \text{uhtc} \xrightarrow{h} 3BB75D$$
8. In this case, the plain-text we wanted is “spkn”.

The location of $y=h(x)$

Chain: $t_0 \xrightarrow{h} h_0 \xrightarrow{r_0} t_1 \xrightarrow{h} h_1 \xrightarrow{r_1} t_2 \xrightarrow{h} h_2$

$y=h(x)$

$y = h_2$
 $x = t_2$

$y=h(x)$

$h(r_1(y)) = h_2$
 $x = t_1$

$y=h(x)$

$h(r_1(h(r_0(y)))) = h_2$
 $x = t_0$





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Tests

- Chain length;



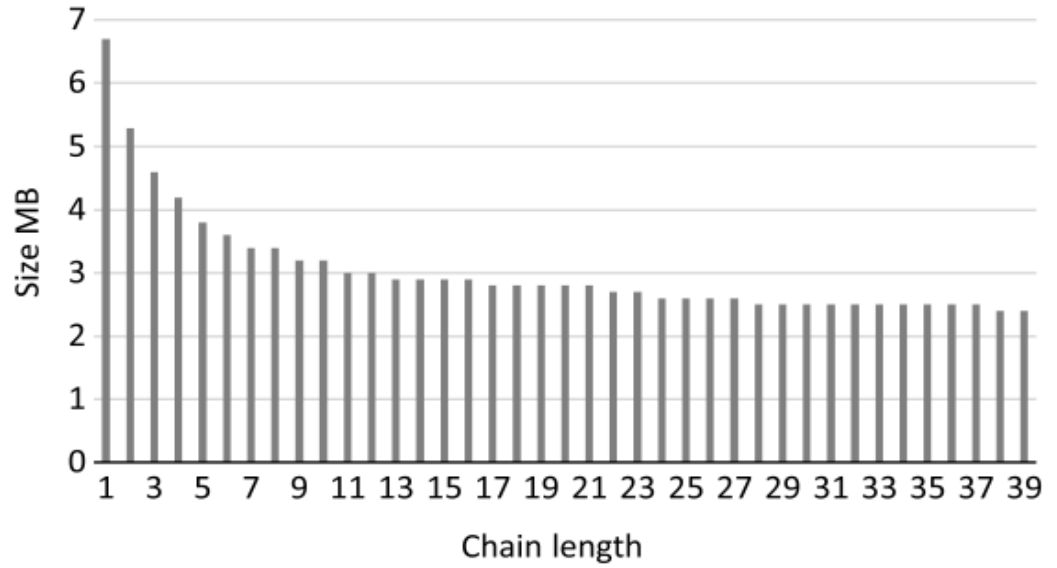


Fig. 2. The size of Rainbow Tables in dependence on the length of chain.

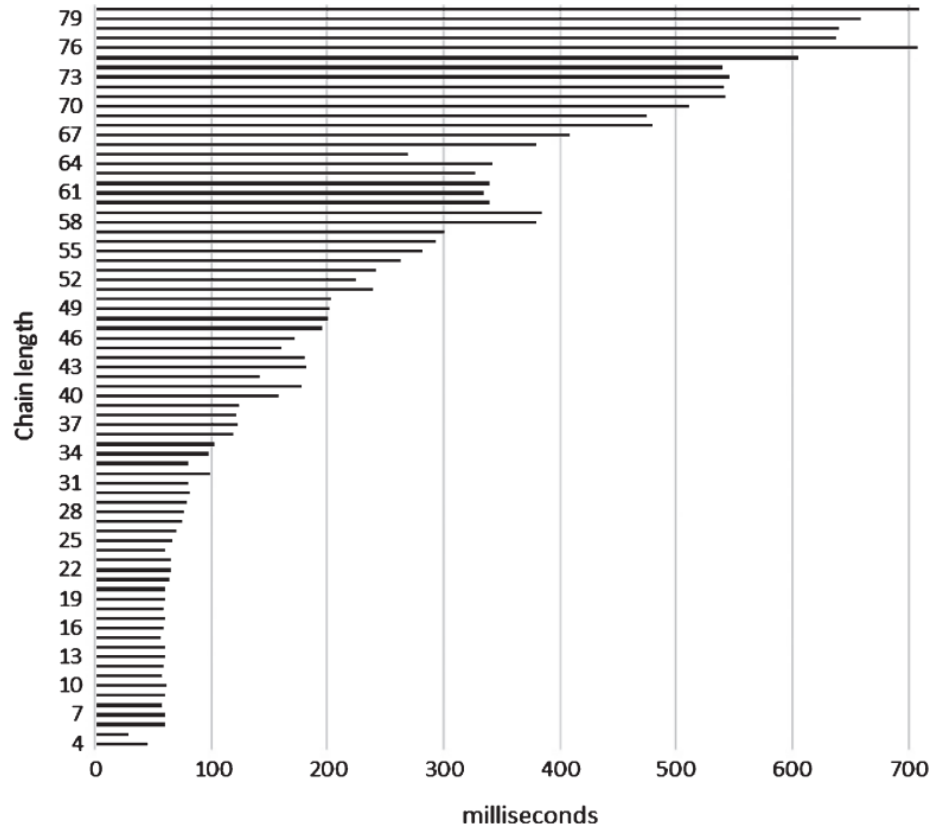


Fig. 1. Dependence on the chain length of the average time to find the password.





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Conclusion

- Time & space;





Acknowledgement

- Great thanks to Peter Dolan, KK Lamberty, and Elena Machkasova.



Reference

1. Horáleka, J., Holík, F., Horák, O., Petr, L., & Sobeslav, V. (2017). Analysis of the use of Rainbow Tables to break hash. *Journal of Intelligent & Fuzzy Systems*,32(2):1523 – 1534, 2017
2. G. Avoine and X. Carpent. Heterogeneous rainbowtable widths provide faster cryptanalyses. In *Proceedings of the 2017 ACM on Asia Conference on Computer and Communications Security, ASIA CCS'17*, pages 815–822, New York, NY, USA, 2017. ACM.
3. Cryptographic hash function. Cryptographic hashfunction — Wikipedia, the free encyclopedia, 2019.[Online; accessed 15-March-2019]

Q&A

