# Enhancements of Security in Wireless Sensor Networks

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#### What are Wireless Sensor Networks?

Wireless sensor networks (WSNs)

- Large number of sensor nodes that have computation and communication capabilities.
- Nodes gather and send the data in various settings
- Use the data locally or allows accessed by some user in real-time.



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• Since WSNs consist of such a **large network** of interconnected components thus making it hard to keep each node secure.

#### • WSNs face resource restrictions:

- storage capacity
- processing power
- battery power
- Security Requirements: Data integrity, Confidentiality and Availability

# Outline



- 4 Phase AA (Anonymous Authentication) Scheme
- 3 Security Analysis

#### 4 Conclusions

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Image: A math

# Outline

#### Background

- Common Attacks on WSNs
- Cryptography used in WSNs
- 2 4 Phase AA (Anonymous Authentication) Scheme
- 3 Security Analysis
- 4 Conclusions

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# Impersonation Attack

The goal is to access unauthorized information acting as a certain party.



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## Node Capture Attack

Goal is to get control of any number of nodes, isn't difficult for them to add, remove, or alter information about some node.



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# Cryptography used in WSNs

Private Key Cryptography

- Goal is to send messages between parties without any adversary listening
- Uses a shared secret key that allows both parties to encrypt and decrypt a message

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**Key:** Data that is used to encrypt and decrypt information. **Encryption** and **Decryption:** functions that take in a message and key as parameters and output specific message



Each letter is equal to its third letter down the alphabet  $(a \rightarrow d, b \rightarrow e, c \rightarrow f)$ 

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- Denoted by h(.) where . are various parameters
- Hash Function: function that can be used to map data of arbitrary size to data of fixed size.
- "One way": Easy to compute in one direction but hard to compute in the other direction.

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- Hash Function: function that can be used to map data of arbitrary size to data of fixed size.
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• Returning hash value from function is used to verify some type of information.

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Background



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# XOR (Exclusive-OR Operator)

- Denoted by
- Means "one or the other but NOT both"
- Example: 11010 ⊕ 10110 = 01100

#### **XOR Truth Table**

а	b	a⊕b
0	1	1
1	1	0
0	0	0
1	0	1

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# XOR (Exclusive-OR Operator)

- Denoted by  $\oplus$
- Means "one or the other but NOT both"
- Example: 11010 ⊕ 10110 = 01100
- Property:

 $\textbf{k} \oplus \textbf{k} \oplus \textbf{m} = \textbf{m} \oplus \ \textbf{k} \oplus \textbf{k} = \textbf{m}$ 

#### **XOR Truth Table**

а	b	a⊕b
0	1	1
1	1	0
0	0	0
1	0	1

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#### Concatenation

- Joining multiple given elements
- Denoted by ||
- Example: 11 || 00 = 1100

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# Outline



#### 4 Phase AA (Anonymous Authentication) Scheme

3 Security Analysis

#### 4 Conclusions

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- Protocol proposed by Gope et al
- Two factor authentication: requires more than just username and password for authentication but also some information that only the user has/knows such as some **physical token**

What does this protocol do?

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What does this protocol do?

• Provides a secure way for user to register with and authenticate with desired sensor node in network.

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What does this protocol do?

- Provides a secure way for user to register with and authenticate with desired sensor node in network.
- Also provides secure way of changing password securely and introducing new new nodes into the network.

Phases:

- Phase 1: Registration Phase
- Phase 2: Anonymous Authentication and Key Exchange
- Phase 3: Password Renewal Phase
- Phase 4: Dynamic Node Addition Phase

## Phase 1: Registration Phase

 Purpose: sets up the underlying foundation for the construction of the proposed authentication scheme.

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# Phase 1: Registration Phase



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# Phase 2: Anonymous Authentication and Key Exchange

• Purpose: Achieve secure authentication among the user, base station, and the sensor node.



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Image: A math



**O Generate:** random number  $N_{II}$ 

Smart Card has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number ID<sub>U</sub>: User identity



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- **Generate:** random number N<sub>U</sub>
- 2 Derive:
  - One Time Alias Identity  $AID_U = h(ID_U||K_{ub}||N_U||S)$

Smart Card has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity



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• 
$$N_x = K_{ub} \oplus N_U$$

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$$V_1 = h(AID_U||K_{ub}||N_x||S)$$

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- Oreate M1

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- Oreate M1
  - M1: {*AID*<sub>U</sub>,*N*<sub>x</sub>, *V*<sub>1</sub>,*S*,*TS*<sub>ub</sub>}

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Upon receiving M1: { $AID_U$ ,  $N_x$ ,  $V_1$ , S,  $TS_{ub}$ } B:

• Checks: *TS*<sub>ub</sub> valid?

Base station has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_U$ : One-time alias identity



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- **Observe 2 Decrypt:**  $ID_U$  and shared key  $K_{ub}$  with validated  $TS_{ub}$  and private key  $\omega$

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• Now *B* knows who the user is.

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#### 4 Phase AA Scheme

# Phase 2: Step 1

Upon receiving M1: { $AID_U$ ,  $N_x$ ,  $V_1$ , S,  $TS_{ub}$ } B:

#### • Checks: *TS*<sub>ub</sub> valid?

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  - Now *B* knows who the user is.
- Oheck if:
  - $V_1 = h(AID_U||K_{ub}||N_x||S)$

Base station has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_U$ : One-time alias identity



#### 4 Phase AA Scheme

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Oheck if:

• 
$$V_1 = h(AID_U||K_{ub}||N_x||S)$$

Base station has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_U$ : One-time alias identity



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Oheck if:

- $V_1 = h(AID_U||K_{ub}||N_x||S)$
- $AID_U = h(ID_U||K_{ub}||N_U||S)$

Base station has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_{U1}$ : One-time alias identity



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#### Generate:

• Session Key SK (randomly)

Base station has:  $K_{ub}$ : Shared key between base station and sensor node  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_{IJ}$ : One-time alias identity



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#### Generate:

- Session Key SK (randomly)
- Time Stamp T

Base station has:  $K_{ub}$ : Shared key between base station and sensor node  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_{IJ}$ : One-time alias identity



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#### Generate:

- Session Key SK (randomly)
- Time Stamp T
- 2 Encrypt: SK such that:
  - $SK' = h(K_{bs}) \oplus SK$

Base station has:  $K_{ub}$ : Shared key between base station and sensor node  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_{II}$ : One-time alias identity



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#### Generate:

- Session Key SK (randomly)
- Time Stamp T
- 2 Encrypt: SK such that:
  - $SK' = h(K_{bs}) \oplus SK$
- Compute:
  - $V_2 = h(AID_U||SK'||T||K_{bs})$

Base station has:  $K_{ub}$ : Shared key between base station and sensor node  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_U$ : One-time alias identity



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#### Generate:

- Session Key SK (randomly)
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- 2 Encrypt: SK such that:
  - $SK' = h(K_{bs}) \oplus SK$
- Compute:
  - $V_2 = h(AID_U||SK'||T||K_{bs})$
- Oreate M2:

Base station has:  $K_{ub}$ : Shared key between base station and sensor node  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_U$ : One-time alias identity



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#### Generate:

- Session Key SK (randomly)
- Time Stamp T
- 2 Encrypt: SK such that:
  - $SK' = h(K_{bs}) \oplus SK$
- Compute:
  - $V_2 = h(AID_U||SK'||T||K_{bs})$
- Create M2:
  - M2: {AID<sub>U</sub>,SK', T,K<sub>bs</sub>}

Base station has:  $K_{ub}$ : Shared key between base station and sensor node  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $ID_B$  Base station identity  $V_1$ : Validation message  $AID_U$ : One-time alias identity



Sensor node receives:  $M2: \{AID_U, SK', T, K_{bs}\}$ 

• Verify if:

• T Valid

Sensor node has:  $K_{ub}$ : Shared key between user and base station SK': Encrypted session key  $ID_U$ : User identity  $ID_B$  Base station identity  $V_2$ : Validation message  $AID_{II}$ : One-time alias identity



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Sensor node receives: M2: {AID<sub>U</sub>,SK', T,K<sub>bs</sub>}

- Verify if:
  - T Valid
  - V<sub>2</sub> =

Sensor node has:  $K_{ub}$ : Shared key between user and base station SK': Encrypted session key  $ID_U$ : User identity  $ID_B$  Base station identity  $V_2$ : Validation message  $AID_U$ : One-time alias identity



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Sensor node receives: M2: {AID<sub>U</sub>,SK', T,K<sub>bs</sub>}

- **O** Verify if:
  - T Valid
  - $V_2 = h(AID_U||SK'||T||V_2)$

Sensor node has:  $K_{ub}$ : Shared key between user and base station SK': Encrypted session key  $ID_U$ : User identity  $ID_B$  Base station identity  $V_2$ : Validation message  $AID_U$ : One-time alias identity



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Sensor node receives:  $M2: \{AID_U, SK', T, K_{bs}\}$ 

- Verify if:
  - T Valid
  - $V_2 = h(AID_U||SK'||T||V_2)$
- Observe the second s

Sensor node has:  $K_{ub}$ : Shared key between user and base station SK': Encrypted session key  $ID_U$ : User identity  $ID_B$  Base station identity  $V_2$ : Validation message  $AID_{II}$ : One-time alias identity



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Sensor node receives:  $M2: \{AID_U, SK', T, K_{bs}\}$ 

# Verify if:

- T Valid
- $V_2 = h(AID_U||SK'||T||V_2)$
- Observe the second s
  - $SK = h(K_{bs}) \oplus SK'$

Sensor node has:  $K_{ub}$ : Shared key between user and base station SK': Encrypted session key  $ID_U$ : User identity  $ID_B$  Base station identity  $V_2$ : Validation message  $AID_U$ : One-time alias identity





Generate: New timestamp T'

Sensor node has: K<sub>bs</sub>: Shared key between base station and sensor SK': Encrypted session key S: Sensor node identity AID<sub>U</sub>: One-time alias identity



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#### Generate: New timestamp T'

- Ompute:
  - $V_3 = h(SK'||K_{bs}||S||T')$

Sensor node has:  $K_{bs}$ : Shared key between base station and sensor SK': Encrypted session key S: Sensor node identity  $AID_U$ : One-time alias identity



- Generate: New timestamp T'
- Ompute:
  - $V_3 = h(SK'||K_{bs}||S||T')$
- Oreate M3:

Sensor node has:  $K_{bs}$ : Shared key between base station and sensor SK': Encrypted session key S: Sensor node identity  $AID_U$ : One-time alias identity



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- Generate: New timestamp T'
- Ompute:
  - $V_3 = h(SK'||K_{bs}||S||T')$
- Oreate M3:
  - *M*3: {*T*′,*S*,*V*<sub>3</sub>}

Sensor node has:  $K_{bs}$ : Shared key between base station and sensor SK': Encrypted session key S: Sensor node identity  $AID_U$ : One-time alias identity



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B: receives  $M3 : \{T', S, V_3\}$ 



Base station has:  $K_{bs}$ : Shared key between base station and sensor  $K_{ub}$ : Shared key between user and base station  $N_{U}$ : Randomly generated number by user  $ID_{U}$ : User identity



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- *B*: receives *M*3 : {*T*′,*S*,*V*<sub>3</sub>}
  - Check: T'
  - **2** Verify if:  $V_3 =$

Base station has:  $K_{bs}$ : Shared key between base station and sensor  $K_{ub}$ : Shared key between user and base station  $N_{U}$ : Randomly generated number by user  $ID_{U}$ : User identity



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*B*: receives *M*3 : {*T*′,*S*,*V*<sub>3</sub>}

- Check: T'
- **2** Verify if:  $V_3 = h(SK'||K_{bs}||S||T')$

Base station has:  $K_{bs}$ : Shared key between base station and sensor  $K_{ub}$ : Shared key between user and base station  $N_U$ : Randomly generated number by user  $ID_{U}$ : User identity



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- *B*: receives *M*3 : {*T*′,*S*,*V*<sub>3</sub>}
  - Check: T'
  - **2** Verify if:  $V_3 = h(SK'||K_{bs}||S||T')$
  - Incrypt: 3
    - TSub so that

$$TS = h(K_{ub}||ID_U||N_U) \oplus TS_{ub}$$

Base station has:  $K_{bs}$ : Shared key between base station and sensor  $K_{ub}$ : Shared key between user and base station  $N_U$ : Randomly generated number by user  $ID_{II}$ : User identity



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- *B*: receives *M*3 : {*T*′,*S*,*V*<sub>3</sub>}
  - Check: T'
  - **2** Verify if:  $V_3 = h(SK'||K_{bs}||S||T')$
  - Incrypt:
    - TSub so that

$$TS = h(K_{ub}||ID_U||N_U) \oplus TS_{ub}$$
  
•  $SK'' = h(K_{ub}||ID_U||N_U) \oplus TS_{ub}$ 

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#### Compute:

•  $V_4 = h(SK''||K_{ub}||N_U||TS)$ 

Base station has:  $K_{ub}$ : Shared key between user and base station SK'': Encrypted session key  $ID_U$ : User identity TS: Encrypted transaction number  $N_U$ : Randomly generated number

by user



#### Compute:

- $V_4 = h(SK''||K_{ub}||N_U||TS)$
- Oreate M4:

Base station has:  $K_{ub}$ : Shared key between user and base station SK'': Encrypted session key  $ID_U$ : User identity TS: Encrypted transaction number  $N_U$ : Randomly generated number by user



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#### Compute:

- $V_4 = h(SK''||K_{ub}||N_U||TS)$
- 2 Create M4:
  - *M*4: {*SK*″,*TS*,*V*<sub>4</sub>}

Base station has:  $K_{ub}$ : Shared key between user and base station SK'': Encrypted session key  $ID_U$ : User identity TS: Encrypted transaction number  $N_U$ : Randomly generated number by user



SC recieves M4: {SK", TS, V<sub>4</sub>}

• Verify:  $V_4 =$ 

Smart Card has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $N_U$ : Randomly generated number by user



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SC recieves M4:  $\{SK'', TS, V_4\}$ 

• Verify:  $V_4 = h(SK''||K_{bs}||N_U||TS)$ ??

Smart Card has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $N_U$ : Randomly generated number by user



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SC recieves M4:  $\{SK'', TS, V_4\}$ 

- Verify:  $V_4 = h(SK''||K_{bs}||N_U||TS)$ ?
- **Decrypt:** *TS*

• 
$$TS_{ub} = h(K_{ub}||ID_U||N_U) \oplus TS$$

Smart Card has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $N_U$ : Randomly generated number by user



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SC recieves M4: {SK", TS, V<sub>4</sub>}

• Verify:  $V_4 = h(SK''||K_{bs}||N_U||TS)$ ?

#### Output: TS

- $TS_{ub} = h(K_{ub}||ID_U||N_U) \oplus TS$
- SC stores TS<sub>ub</sub> in its memory for future authentication.

Smart Card has:  $K_{ub}$ : Shared key between user and base station  $TS_{ub}$ : Transaction number  $ID_U$ : User identity  $N_U$ : Randomly generated number by user



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SC recieves M4: {SK", TS, V<sub>4</sub>}

- Verify:  $V_4 = h(SK''||K_{bs}||N_U||TS)$ ??
- Output: TS
  - $TS_{ub} = h(K_{ub}||ID_U||N_U) \oplus TS$
  - SC stores TS<sub>ub</sub> in its memory for future authentication.

User is now authenticated and free to communicate with said node *S*.

## Phase 3 and Phase 4

Phase 3: Password Renewel Phase

- Purpose: Provide secure manner for user to change password on smart card
- Unlike most AA schemes the user need not communicate with *B* and is free to change his/her password on the smart card *SC*

Phase 4: Dynamic Node Addition Phase

Purpose: Provide a secure manner of adding new nodes into the network
## Outline

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#### Background

#### 2 4 Phase AA (Anonymous Authentication) Scheme

#### Security Analysis

- Resilience Against Key Compromise Impersonation Attack
- Resilience Against Node Capturing Attack

#### 4) Conclusions

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## Resilience Against Key Compromise Impersonation Attack

Suppose an adversary *A* has obtained the servers' (*B*) secret key  $\omega$  and the encoded parameters from the servers database.



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#### **Resilience Against Node Capturing Attack**

How do captured nodes affect entire network?

Suppose a compromised node has acquired the private key  $\omega$  and session key *SK*.

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#### **Resilience Against Node Capturing Attack**



## Outline

#### Background

- 2 4 Phase AA (Anonymous Authentication) Scheme
- 3 Security Analysis
- 4 Conclusions

#### How does this protocol enhance security in WSNs?

• Data Integrity: One of the main ways data remains unaltered through transmission is defense against node capture attack.

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### How does this protocol enhance security in WSNs?

- Data Integrity: One of the main ways data remains unaltered through transmission is defense against node capture attack.
- Confidentiality: Since the protocol defends against an impersonation of the user then we can say their identity will stay secure.
- Availability: Once the user is securely authenticated through our protocol any authorized user can acquire their data in real-time.

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Conclusions

#### Acknowledgments

# Thank you to Nic McPhee and Elena Machkasova for feedback and help during process. As well as peers that gave me feedback.

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Conclusions

Thanks for coming!

## **Questions?**

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